

RLH: Bitmap Compression Technique Based on Run-Length and Huffman Encoding

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Presentation Outline

- **⇒** Bitmap index concept and main characteristics
- Reducing size of bitmap index
- Run Length Huffman compression algorithm
- Run Length Huffman experimental evaluation
- Summary



Bitmap Index

- Composed of bitmaps
- **○** A bitmap is a vector of bits
 - every value from a domain has its own bitmap
 - the number of bits = the number of records
 - a given bit corresponds to a given record
- Basic characteristics
 - Efficient in answering equality and range queries
 - BI size depends on the cardinality of an indexed attribute
 - large for high cardinality attributes

ID sex	
1 male	
2 fema.	le
3 fema.	
4 fema.	le
5 male	
6 male	
7 male	
8 fema.	
9 fema.	le
10 male	
11 male	
12 male	
13 fema.	
14 femal	
15 fema.	le
16 male	
17 femal	
18 femal	
19 femal	le

Clients

bitmap	index
$_{ m female}$	$_{ m male}$
0	1
1	0
1	0
1	0
0	1
0	1
0	1 0
1	
1	0
0	1
0	1
0	1
1	0
1	0
1	0
0	1
1	0
1	0
1	0



Reducing BI Size

- Binning
 - [Kou00, SWS04, RSW05]
- Encoding
 - [WuBu98, ChIo99]
- Compressing
 - Byte-aligned Bitmap Compression [AnZi96]
 - Word-Aligned Hybrid [SWS02, WOS04]
 - Approximate Encoding [ACFT06]
 - may create false positives, additional verification
 - Reordering [JKCKV04, PTF05]
 - computationally very complex
 - reordering heuristics



Bitmap Compression (1)

- **⇒** Byte-aligned Bitmap Compression (BBC)
- Word-Aligned Hybrid (WAH)
- Based on the run-length encoding
 - homogeneous vectors of bits are replaced with a bit value (0 or 1) and the vector length
 - **■** 0000000 1111111111 000 **⇒** 07 110 03
- **⇒** BBC and WAH
 - a bitmap is divided into words
 - BBC uses 8-bit words
 - WAH uses 31-bit words



Bitmap Compression (2)

- **WAH-compressed bitmaps are larger than BBC-compressed ones**
- ◆ Operations on WAH-compressed bitmaps are faster than on BBC-compressed ones [SWS02, WOS02, WOS04]
- Our further focus is on comparing WAH to our approach



WAH (1)

a) an example bitmap being compressed (5456 bits) 5394 bits having value "0" 31 bits 31 bits b) dividing the bitmap into 31-bits groups 31 bits 31 bits 31 bits group 1 group 2 group 176 c) merging adjacent homogeneous groups 31 bits 174 * 31 bits 31 bits group 1 group 2-175 group 176 d) group encoding by means of a 32-bits word **0** 100000.....0001110000111 **1 0** .000...0010101110. **0** 00111111111.....1111011111, fill length 174 * 31 bits 31 bits of the last group 31 bits of the first group bit=0: fill value bit=0: tail word bit=0: tail word bit=1: fill word

run 2

Example taken from [StWu07]

run 1



WAH (2)

- 1. For low cardinality attributes bitmaps are dense
 - many homogeneous 31-bit words filled with 1
- 2. For high cardinality attributes bitmaps are sparse
 - many homogeneous 31-bit words filled with 0
- 3. For medium cardinality attributes
 - the number of homogeneous 31-bit words is lower
 - the compression ratio decreases
- **○** A need for bitmap compression technique suitable for medium cardinality attributes



Our Approach: RLH

- RLH the Run-Length Huffman Compression
- Based on
 - the Huffman encoding
 - a modified run-length encoding



Huffman Encoding

Concept

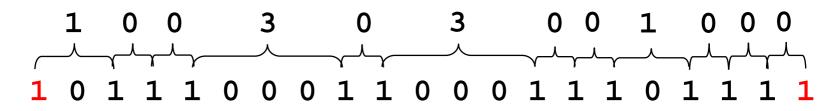
- original symbols from a compressed file are replaced with bit strings
- the more frequently a given symbol appears in the compressed file the shorter bit string for representing the symbol
- encoded symbols and their corresponding bit strings are represented as a Huffman tree
- the Huffman tree is used for both compressing and decompressing



RLH (1)

Modified run-length encoding

measures and encoded distances between bits of value 1



Cl	ients

ID	sex	$_{ m female}$	Γ
1	male	0	
2	female	1	Γ
3	female	1	
4	female	1	
5	male	0	
6	male	0	Γ
7	$_{ m male}$	0	Γ
8	female	1	Γ
9	female	1	
10	male	0	
11	male	0	Γ
12	male	0	
13	female	1	
14	female	1	
15	female	1	
16	$_{ m male}$	0	
17	female	1	
18	female	1	Γ

female

bitmap index female | male

0

0

female: 100303001000

male: 030020033

⇒ Bitmaps encoded this way are input for the Huffman encoding

LAP 2007, Portugal



RLH (2)

Huffman encoding

step1: computing frequencies of symbols (distances) in encoded bitmaps

female: 100303001000

male: 030020033



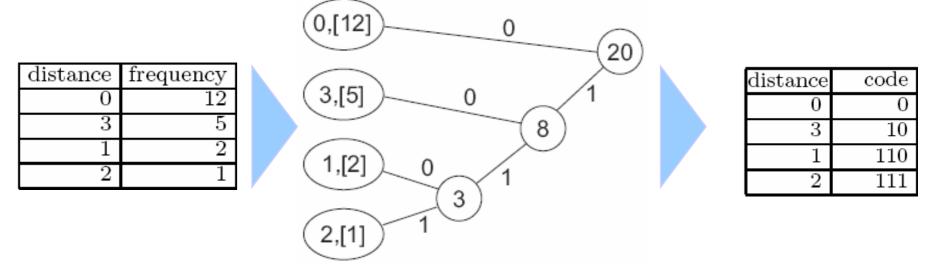
distance	frequency
0	12
3	5
1	2
2	1



RLH (3)

Huffman encoding

step2: building a Huffman tree



 an encoded symbol is represented by a path from the root to a leaf



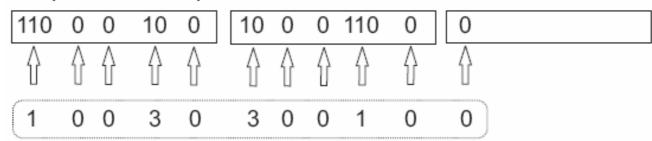
RLH (4)

Huffman encoding

step3: replacing distances with their Huffman codes

$\operatorname{distance}$	code
0	0
3	10
1	110
2	111

compressed bitmap for sex='female'



the result of modified run-length encoding for bitmap sex='female'

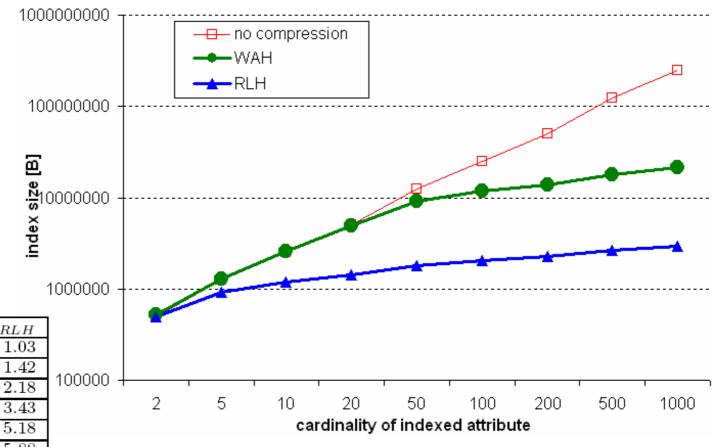


Experimental Evaluation

- Comparing RLH, WAH, and uncompressed bitmaps with respect to
 - bitmap sizes
 - query response times
- **⇒** Implementation in Java
 - data and bitmap indexes stored on disk in OS files
- **⇒** Experiments run on
 - PC, AMD Athlon XP 2500+; 768 MB RAM; Windows XP
- **Data**
 - 2 000 000 indexed rows
 - indexed attribute of type integer
 - cardinality from 2 to 1000
 - randomly distributed values



WAH and RLH: index sizes

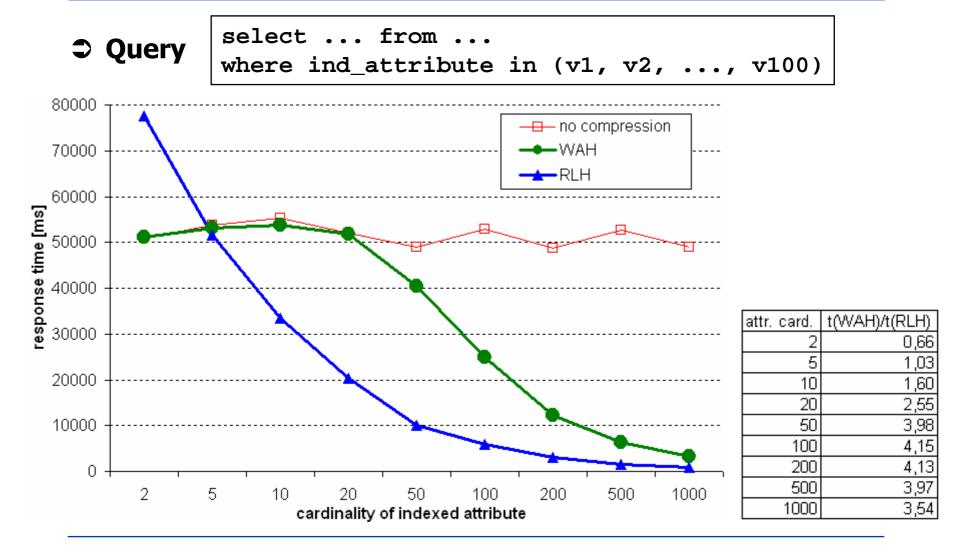


attr.cardinality	s_{WAH}/s_{RLH}
2	1.03
5	1.42
10	2.18
20	3.43
50	5.18
100	5.88
200	6.01
500	6.68
1000	7.36

cardinality increases
 size of RLH bitmaps increases more slowly than WAH



WAH and RLH: response times





Updating RLH Bitmaps

Costly process

- decompressing the whole bitmap
- modifying the bitmap
- compressing the bitmap

Updating a RLH bitmap

- changes frequencies of distances between 1 bits
- creates new distances between 1 bits
- requires building a new Huffman tree

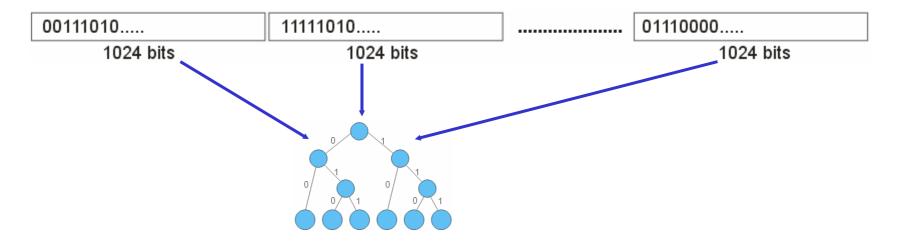
⇒ In a DW environment index structures

- are dropped before loading a DW
- are recreated after loading is finished



RLH1024 Compression (1)

- 1. Dividing a bitmap into 1024-bit sections
 - constructing one Huffman tree based on frequencies of distances from all 1024-bit sections



- 2. Including in the HT all possible distances that may appear in a 1024-bit section
 - non-existing distances have assigned the frequency of 1



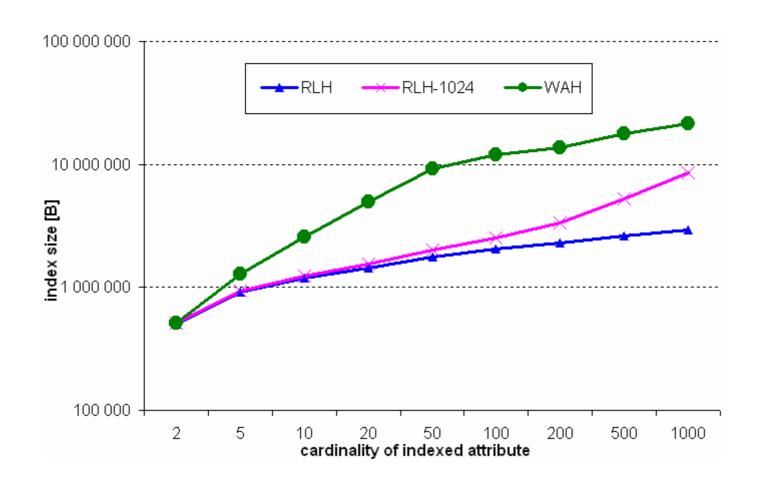
RLH1024 Compression (2)

Advantages

- including all the possible distances in the HT eliminates the need of rebuilding the HT after such a bitmap update that results in a new distance
- 1024-bit sections can be read and processed in parallel
- in order to update a bitmap, only an appropriate
 1024-bit section has to be read and uncompressed

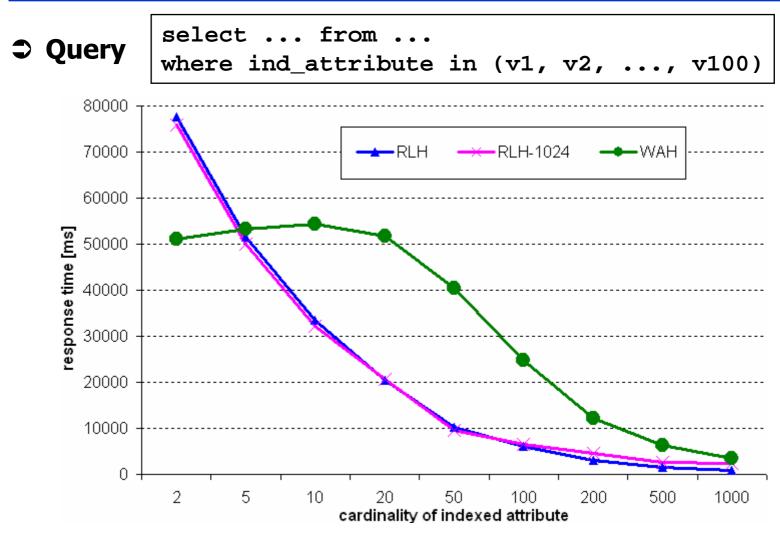


RLH and RLH1024: index sizes





RLH and RLH1024: response times





Summary

- Alternative bitmap compression technique based on the run-length encoding and on Huffman encoding
 - RLH
 - RLH1024
- Observations
 - RLH offers a higher efficiency in accessing data than WAH for attribute cardinality from 5 to 1000
 - Bitmaps compressed with RLH are much smaller than corresponding bitmaps compressed with WAH for attribute cardinalities >10
 - RLH1024 offers a data access time characteristic similar to RLH, but additionally RLH1024 may better support bitmap updates



Ongoing and Future Work

Ongoing

- evaluating the impact of values distribution on WAH and RLH
- evaluating other than 1024-bit partition schemes
- evaluating the efficiency of updating bitmaps in RLH and RLH1024

⇒ Future

- developing a cost model for RLH
- developing a framework for selecting the most efficient bitmap partition scheme for RLH
- developing a framework for selecting the most efficient bitmap compression technique for a given data characteristic
- experimentally comparing BBC, WAH, RLH, and AE
- integrating RLH into FastBit



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Binning

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BBC

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